Why is modal logic decidable

CS 112 Fall 2018

Slides by Petros Potikas

Computational problems

Two computational problems:

1. *Model-checking* problem: is a given formula true at a given state at a given Kripke structure
2. *Validity* problem: is a given formula true in all states of all Kripke structures

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- But in ML we have arbitrary nesting of modalities.
- So, this cannot be captured by bounded quantifier alternation.

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- FP^2 does not enjoy the nice computational properties of FO^2.
- Decidability of CTL can be explained by *tree-model property*, which is enjoyed by CTL, but not by FP^2.
- Finally, the tree model property leads to automata-based decision procedures.

Syntax

Definition

(The Basic Modal Language) Let $\mathbb{P} = \{\mathbb{P}_0, \mathbb{P}_1, \mathbb{P}_2, ...\}$ be a set of sentence letters, or atomic propositions. We also include two special propositions $\top$ and $\bot$ meaning 'true' and 'false' respectively. The set of well-formed formulas of modal logic is the smallest set generated by the following grammar: $\mathbb{P}_0, \mathbb{P}_1, \mathbb{P}_2, ... \mid \top \mid \bot \mid \neg A \mid A \vee B \mid A \wedge B \mid A \rightarrow B \mid \Box A \mid \Diamond A$

Examples

Modal formulas include: $\Box\bot, \mathbb{P}_0 \rightarrow \Diamond(\mathbb{P}_1 \wedge \mathbb{P}_2)$.

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Truth conditions:

1. $(M, s) \models \mathbb{P}_i$ iff $s \in \pi(\mathbb{P}_i)$
2. $(M, s) \models \top$
3. $(M, s) \not\models \bot$
4. $(M, s) \models \neg A$ iff not $(M, s) \models A$
5. $(M, s) \models A \vee B$ iff either $(M, s) \models A$ or, $(M, s) \models B$,or both
6. $(M, s) \models \Box A$ iff for every t, s.t. $R(s, t), (M, t) \models A$

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6. $(M, s) \models \Box A$ iff for every t, s.t. $R(s, t), (M, t) \models A$

- A sentence true at every possible world in every model is said to be *valid*, written $\models A$

Model-checking problem

Theorem

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Proof.

Let $\phi_1, ..., \phi_m$ be the subformulas of ϕ listed in order of length. Thus $\phi_m = \phi$, and if ϕ_i is a subformulas of ϕ_j, then $i < j$. There are at most $|\phi|$ subformulas, so $m \leq |\phi|$. By induction on k, we can show that we can label each state s with ϕ_j or $\neg\phi_j$, for $j = 1, ..., k$, depending on whether or not ϕ_j is true in s in time $O(k||M||)$. Only interesting case is $\phi_{k+1} = \Box\phi_j$, $j < k + 1$. By induction hypothesis, we have that each state has already been labeled with ϕ_j or $\neg\phi_j$, so we know if node s can be labeled with ϕ_{k+1} or not in time $O(||M|||)$. □

Characterizing the properties of necessity

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Two approaches:

1. *Proof-theoretic*: all properties of necessity can be formally derived from a short list of basic properties
2. *Algorithmic*: we study algorithms that recognize properties of necessity and consider their computational complexity.

Properties of necessity

Some basic properties of necessity:

Theorem

For all formulas ϕ, ψ, and Kripke structures M:

1. *if ϕ is an instance of a propositional tautology, then $M \models \phi$*
2. *if $M \models \phi$ and $M \models \phi \rightarrow \psi$, then $M \models \psi$*
3. $M \models (\Box\phi \wedge \Box(\phi \rightarrow \psi)) \rightarrow \Box\psi$
4. *if $M \models \phi$, then $M \models \Box\phi$*

Characterizing the properties of necessity: Proof-theoretic

Consider the following axiom system $\mathcal{K}$:

- (A1) All tautologies of propositional calculus
- (A2) $(\Box\phi \wedge \Box(\phi \rightarrow \psi)) \rightarrow \Box\psi$ (Distribution axiom)
- (R1) From ϕ and $\phi \rightarrow \psi$ infer ψ (Modus ponens)
- (R2) From ϕ infer $\Box\phi$ (Generalization)

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Theorem (Kripke '63)

$\mathcal{K}$ is a sound and complete axiom system.

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Theorem (Fischer,Ladner '79)

If a modal formula ϕ is satisfiable, then ϕ is satisfiable in a Kripke structure with at most $2^{|\phi|}$ states.

Characterizing the properties of necessity: algorithmically

- From the above Theorem we can get an algorithm (not efficient) for testing validity of a formula ϕ: construct all Kripke structures with at most $2^{|\phi|}$ states and check if the formula is true in every state of each of these structures.

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- From the above Theorem we can get an algorithm (not efficient) for testing validity of a formula ϕ: construct all Kripke structures with at most $2^{|\phi|}$ states and check if the formula is true in every state of each of these structures.
- The "inherent difficulty" of the problem is given by the next theorem:

Theorem (Ladner '77)

The validity problem for modal logic is PSPACE-complete.

Modal logic vs. First-Order Logic

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- Every Kripke structure M can be viewed as a relational structure M^* over the vocabulary $\mathbb{P}^*$.
- Formally, a mapping from a Kriple structure $M = (S, \pi, R)$ to a relational structure M^* over the vocabulary $\mathbb{P}^*$ has:
  1. domain of M^* is S.
  2. for each propositional constant $q \in \mathbb{P}$, the interpretation of q in M^* is the set $\pi(q)$.
  3. the interpretation of the binary predicate $\mathcal{R}$, is the binary relation R.

Translation of Modal logic to First-Order Logic

A translation from modal formulas into first-order formulas over the vocabulary $\mathbb{P}^*$, so that for every modal formula ϕ there is corresponding first-order formula ϕ^* with one free variable (ranging over S):

1. $q^* = q(x)$ for a propositional constant q
2. $(\neg\phi)^* = \neg(\phi^*)$
3. $(\phi \wedge \psi)^* = (\phi^* \wedge \psi^*)$
4. $(\Box\phi)^* = (\forall y(R(x, y) \rightarrow \phi^*(x/y)))$, where y is a new variable not appearing in ϕ^* and $\phi^*(x/y)$ is the result of replacing all free occurrences of x in ϕ^* by y

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Example

$(\Box\Diamond q)^* = \forall y(R(x,y) \rightarrow \exists z(R(y,z) \wedge q(z)))$

Theorem (vBenthem '74,'85)

1. *$(M, s) \models \phi$ iff $(M^*, V) \models \phi^*(x)$, for each assignment V s.t. $V(x) = s$.*
2. *ϕ is a valid modal formula iff ϕ^* is a valid first-order formula.*

ϕ^* is true of exactly the domain elements corresponding to states s for which $(M, s) \models \phi$

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- Carefully examining propositional modal logic, reveals that it is a fragment of *2-variable first-order logic* (FO^2), e.g. $\forall x \forall y (R(x,y) \rightarrow R(y,x))$ is in FO^2, while $\forall x \forall y \forall z (R(x,y) \wedge R(y,z) \rightarrow R(x,z))$ is not in FO^2.

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2. ϕ *is a valid modal formula iff* ϕ^+ *is a valid* FO^2 *formula.*

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- But Mortimer's proof shows bounded-model property.

Theorem

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- Further, the translation of modal logic to FO^2 is linear, so we have Theorem 5.
- Note, however, that the validity problem for FO^2 is hard for co-NEXPTIME (Fürer81) and also complete, while from Theorem 6 modal logic is PSPACE-complete.
- The embedding to FO^2 does not give a satisfactory explananation of the tractability of modal logic.

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The validity problem for modal logic in M_r is PSPACE-complete.

Theorem

A modal formula ϕ is valid in M_r iff the FO^2 $\forall x(R(x,x) \rightarrow \phi^+)$ is valid.

Axiom system S5

What about other properties of necessity?

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1. *$S5$ is sound and complete for M_{rst}.*
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Symmetry can be expressed by FO^2, $\forall x, y(R(x, y) \rightarrow R(y, x))$, while transitivity cannot $\forall x, y, z(R(x, y) \wedge R(y, z) \rightarrow R(x, z))$.

About decidability of modal logic

- The validity in a modal logic is typically decidable. It is very hard to find a modal logic, where validity is undecidable.
- The translation to FO^2 provides a partial explanation why modal logic is decidable.